

Lecture 6: Queueing Theory III

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Disclaimer: *These notes are primarily adapted from expositional texts, including work by Jayatiprasad Medhi, Karl Sigman, and János Sztrik. These notes are not meant to be complete or fully rigorous; some proofs are not given, incomplete, or only outlined, as they are discussed in class.*

Discussion 6.1. *Continuing our discussion from Lecture 5 on priorities (Discussion 5.4), now let's consider a non-preemptive case, where we have an $M/M/1$ queue with two types of customers, but now the arrival of a type 1 customer does not disrupt the service of a type 2 customer. After the server is done serving a type 2 customer, we start serving type 1 customers if there are any. Find the average number of customers and average waiting times of both types.*

6.1 Jackson Networks

Jackson's network model is defined as follows. Assume that customers from one node (queueing system) i proceed to an arbitrary node, and new customers may arrive to a node from outside (say customers arrive to node i according to a Poisson process with rate λ_i). Suppose that there are k nodes, where the i th node ($i = 1, \dots, k$) consists of c_i exponential servers with parameter μ_i (that is, each node contains a $M/M/c$ queueing system). Customers after receiving service at the i th node proceed to the j th node with probability p_{ij} .

Customers at node i depart from the system with probability

$$q_i = 1 - \sum_{j=1}^k p_{ij}.$$

Consider Jackson's general network model with k nodes. The arrivals can be categorized into two groups: the external arrivals (with rate λ_i) and internal arrivals (with rate $\sum_{j=1}^k p_{ji} \lambda_j$). Therefore, the effective arrival rate to node i (or the effective rate of flow through node i) is

$$\alpha_i = \lambda_i + \sum_{j=1}^k p_{ji} \alpha_j, \quad i = 1, 2, \dots, k; \quad (6.1)$$

where these equations are also referred as traffic (flow balance, conservation, etc.) equations.

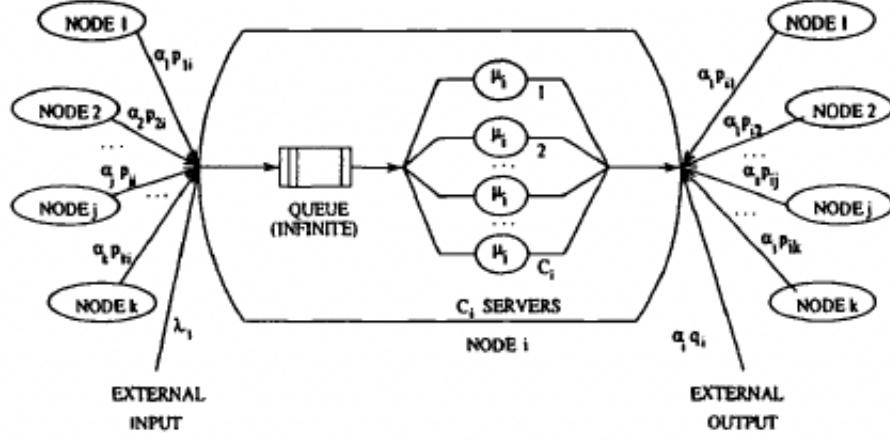


Figure 5.1 Node i in a Jackson Network.

Theorem 6.2 (Jackson's Theorem). Let (n_1, n_2, \dots, n_k) denote the state of the complete system in which there are n_i (in the queue and in service) at node i in a Jackson network of Markovian queues in equilibrium, and let $p(n_1, \dots, n_k)$ be the probability that the system is in the state (n_1, \dots, n_k) .

Assume that

$$\rho_i = \frac{\alpha_i}{\mu_i} < 1, \quad i = 1, 2, \dots, k,$$

where $\{\alpha_i\}$ are given by the balance equations

$$\alpha_i = \lambda_i + \sum_{j=1}^k \alpha_j p_{ji}, \quad i = 1, 2, \dots, k. \quad (6.2)$$

If $p_i(n)$ denotes the probability that there are n customers in the system (in queue plus service) for the $M/M/c_i$ queue with input rate α_i and service rate μ_i for each of the c_i servers, i.e.,

$$p_i(n) = p_i(0) \frac{\left(\frac{\alpha_i}{\mu_i}\right)^n}{n!}, \quad n = 0, 1, 2, \dots, c_i, \quad (6.3)$$

$$= p_i(0) \frac{\left(\frac{\alpha_i}{\mu_i}\right)^n}{c_i! c_i^{n-c_i}}, \quad n = c_i + 1, \dots, \quad (6.4)$$

then we have

$$p(n_1, n_2, \dots, n_k) = p_1(n_1) p_2(n_2) \cdots p_k(n_k). \quad (6.5)$$

Proof Sketch. Let $p_t(n_1, \dots, n_k)$ be the probability that the complete system is in state (n_1, \dots, n_k) at time t .

Let

$$q_i = 1 - \sum_j p_{ij}, \quad a_i(n) = \min\{n_i, c_i\} = \begin{cases} n_i, & \text{if } n < c_i, \\ c_i, & \text{if } n \geq c_i, \end{cases}$$

$$\delta_i = \min\{n_i, 1\} = \begin{cases} 1, & n_i \geq 1, \\ 0, & n_i = 0. \end{cases}$$

Our goal is to write the differential equations satisfied by p_t . Therefore, we will consider the state changes in an infinitesimal interval $(t, t + h)$ following the interval $(0, t)$. p_t . Consider the following four mutually exclusive ways to move from t to $t + h$:

(A) State at t is (n_1, \dots, n_k) and there are no arrivals or departures occur to or from any node externally. We get

$$\Pr(A) = p_t(n_1, \dots, n_k) \left[1 - \left(\sum_i \lambda_i \right) h - \sum_i a_i(n_i) \mu_i h \right] + o(h). \quad (6.6)$$

(B) State at t is $(n_1, \dots, n_i + 1, \dots, n_k)$ and there is one service completion at i in $(t, t + h)$, and this completion departs from the system (with probability q_i). We get

$$\Pr(B) = \sum_{i=1}^k p_t(n_1, \dots, n_i + 1, \dots, n_k) [a_i(n_i + 1) \mu_i q_i] h + o(h). \quad (6.7)$$

(C) State at t is $(n_1, \dots, n_i - 1, \dots, n_k)$ and there is one arrival from the external source to node i in the interval $(t, t + h)$. We get

$$\Pr(C) = \sum_{i=1}^k p_t(n_1, \dots, n_i - 1, \dots, n_k) [\lambda_i h \delta_i] + o(h). \quad (6.8)$$

(D) State at t is $(n_1, \dots, n_i + 1, \dots, n_j - 1, \dots, n_k)$: there is one service completion at node i in $(t, t + h)$, and the one whose service is completed moves to node j with probability p_{ij} . Thus,

$$\Pr(D) = \sum_i \sum_j p_t(n_1, \dots, n_i + 1, \dots, n_j - 1, \dots, n_k) [a_i(n_i + 1) \mu_i h p_{ij}] + o(h). \quad (6.9)$$

Merging all cases, we have

$$p_{t+h}(n_1, \dots, n_k) = \Pr(A) + \Pr(B) + \Pr(C) + \Pr(D). \quad (6.10)$$

(6.10) can be written as

$$p_{t+h}(n_1, \dots, n_k) - \Pr(A) = \Pr(B) + \Pr(C) + \Pr(D). \quad (6.11)$$

$$p_{t+h}(n_1, \dots, n_k) - p_t(n_1, \dots, n_k) \left[1 - \left(\sum_i \lambda_i \right) h - \sum_i a_i(n_i) \mu_i h \right] - o(h) = \Pr(B) + \Pr(C) + \Pr(D). \quad (6.12)$$

Taking the limit as $h \rightarrow 0$ and solving for $p'(t) = 0$ gives the equations satisfied by steady-state probabilities:

$$\begin{aligned} \left[\sum_i \lambda_i + \sum_i a_i(n_i) \mu_i \right] p(n_1, \dots, n_k) &= \sum_i a_i(n_i + 1) \mu_i q_i p(n_1, \dots, n_i + 1, \dots, n_k) \\ &\quad + \sum_i \lambda_i \delta_i p(n_1, \dots, n_i - 1, \dots, n_k) \\ &\quad + \sum_i \sum_j a_i(n_i + 1) \mu_i p_{ij} p(n_1, \dots, n_i + 1, \dots, n_j - 1, \dots, n_k). \end{aligned}$$

Finally, one can show that (6.5) uniquely satisfies the equations above. ■

6.2 The M/G/1 Model

Now assume that we have a Poisson arrival process with rate λ , the service times are i.i.d. and follow a general distribution with mean $\mathbb{E}[S] = \frac{1}{\mu}$, and there is a single server. Again, for stability, we assume that $\rho = \frac{\lambda}{\mu} < 1$.

Let R be the residual service time and let P_k denote the probability that there are k customers in the system in the steady-state. By PASTA property, we have

$$\begin{aligned} W_q &= \sum_{k=1}^{\infty} (E(R) + (k-1)E(S)) P_k \\ &= \sum_{k=1}^{\infty} E(R) P_k + \left(\sum_{k=1}^{\infty} (k-1) P_k \right) E(S) \\ &= E(R)\rho + L_q E(S). \end{aligned}$$

where the equation follows since $1 - \rho = P_0$. By Little's Law, we get

$$W_q = \frac{\rho \mathbb{E}[R]}{1 - \rho} \quad (6.13)$$

which is known as the Pollaczek-Khintchine mean value formula. So now we have to characterize $\mathbb{E}[R]$.

Proposition 6.3. *We have $\mathbb{E}[R] = \frac{\mathbb{E}[S^2]}{2\mathbb{E}[S]}$.*

There are various ways to prove this, but the most pedagogical one is via renewal reward theorem (see the next section). From Proposition 6.3, we have

$$E(R) = \frac{E(S^2)}{2E(S)} = \frac{\text{Var}(S) + E[S]^2}{2E(S)} = \frac{1}{2} (C_S^2 + 1) E(S), \quad (2.33)$$

where C_S^2 is the squared coefficient of the service time S . From here, we get

Thus for the mean waiting time we have

$$W_q = \frac{\rho E(R)}{1 - \rho} = \frac{\rho}{2(1 - \rho)} (C_S^2 + 1) E(S).$$

And by Little's law, we get

$$L_q = \frac{\rho^2}{1 - \rho} \frac{C_S^2 + 1}{2}.$$

Discussion 6.4. *Kingman's G/G/N formula:*

$$W \approx \frac{1}{\mu N} \cdot \frac{\rho \sqrt{2(N+1)-1}}{1 - \rho} \cdot \frac{C_A^2 + C_S^2}{2} + \frac{1}{\mu}.$$

Discussion 6.5. *Matching queues, and the relationship between regret and queue-lengths.*

6.3 The Renewal Reward Theorem

Definition 6.6. *A random point process $\psi = \{t_n\}$ for which the (non-negative) interarrival times $X_n = t_n - t_{n-1}$, $n \geq 1$, form an i.i.d. sequence is called a renewal process.*

Following the definition, t_n is called the n th *renewal epoch* and $F(x) := P(X \leq x)$, $x \geq 0$, denotes the common inter-arrival time distribution. $t_n = X_1 + \dots + X_n$, and $N(t) = \max\{n : t_n \leq t\}$ is the counting process. The rate of the renewal process is denoted by $\lambda \triangleq 1/\mathbb{E}[X]$.

Theorem 6.7 (Elementary renewal theorem). *For a renewal process,*

$$\lim_{t \rightarrow \infty} \frac{N(t)}{t} = \lambda \quad a.s.$$

and

$$\lim_{t \rightarrow \infty} \frac{\mathbb{E}[N(t)]}{t} = \lambda.$$

Proof. We start with the first statement. Note that $t_n = X_1 + \dots + X_n$, $n \geq 1$. Consider the time

$$t_{N(t)} \leq t < t_{N(t)+1}, \quad (6.14)$$

Then we have

$$t_{N(t)} = \sum_{j=1}^{N(t)} X_j, \quad t_{N(t)+1} = \sum_{j=1}^{N(t)+1} X_j,$$

which can be written as

$$\frac{1}{N(t)} \sum_{j=1}^{N(t)} X_j \leq \frac{t}{N(t)} < \frac{1}{N(t)} \sum_{j=1}^{N(t)+1} X_j.$$

By the strong law of large numbers, the left and right hand sides converge to $E(X)$ as $t \rightarrow \infty$ almost surely. One can prove the second part via Wald's identity (which you will prove in your homework with some provided hints). ■

Now let $R(t) = \sum_{j=1}^{N(t)} R_j$ be the total amount of reward collected by time t , where $N(t)$ is the counting process for the renewal process. We want to calculate our long-run reward rate

$$\lim_{t \rightarrow \infty} \frac{R(t)}{t}.$$

Theorem 6.8 (Renewal reward theorem). *For a positive recurrent renewal process in which a reward R_j is earned during cycle length X_j and such that $\{(X_j, R_j) : j \geq 1\}$ is i.i.d. with $\mathbb{E}[|R_j|] < \infty$, the long run rate at which rewards are earned is given by*

$$\lim_{t \rightarrow \infty} \frac{R(t)}{t} = \frac{\mathbb{E}[R]}{\mathbb{E}[X]} = \lambda \mathbb{E}[R] \quad a.s., \quad (6.15)$$

where (X, R) denotes a typical “cycle” (X_j, R_j) ; $\lambda = \{\mathbb{E}[X]\}^{-1}$ is the arrival rate for the renewal process.

Moreover,

$$\lim_{t \rightarrow \infty} \frac{\mathbb{E}[R(t)]}{t} = \frac{\mathbb{E}[R]}{\mathbb{E}[X]}. \quad (6.16)$$

Discussion 6.9. Now let's prove Proposition 6.3 via the renewal reward theorem. Consider a renewal point process $\{t_n : n \geq 1\}$ with i.i.d. interarrival times $X_n = t_n - t_{n-1}$, $n \geq 1$. Define

$$A(t) = t_{N(t)+1} - t, \quad t \geq 0. \quad (6.17)$$

$A(t)$ is called the excess at time t , or remaining lifetime. If $t_{n-1} \leq t < t_n$, then

$$A(t) = t_n - t \leq X_n.$$

Note that if $\{t_n\}$ is a Poisson process at rate λ , then by the memoryless property we have $A(t) \sim \exp(\lambda)$, $t \geq 0$. But for a general renewal process (as in residual service time in $M/G/1$ queue), we need to be smarter. We want to show that

$$\lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t A(s) ds = \frac{\mathbb{E}[X^2]}{2\mathbb{E}[X]} \quad a.s..$$

Note that we can view the i.i.d. X_j as cycle lengths (service times), and $r(t) = A(t)$ as the generated reward rate at time t . Let R_1 be the generated reward in the first cycle. Then we have

$$R_1 = \int_0^{X_1} A(s) ds = \int_0^{X_1} (X_1 - s) ds = \frac{X_1^2}{2}.$$

Since $\{(X_j, R_j)\}$'s are i.i.d., by the renewal reward theorem, we almost surely have

$$\lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t A(s) ds = \frac{\mathbb{E}[R]}{\mathbb{E}[X]} = \frac{\mathbb{E}[X^2]}{2\mathbb{E}[X]}$$

Discussion 6.10 (Inspection paradox). Let $S(t) = t_{N(t)+1} - t_{N(t)}$ be the length of the interarrival time covering time t . If $t_{j-1} \leq t < t_j$, then we have $S(t) = X_j$. Define the reward rate as $r(t) = S(t)$. Then we get

$$R_j = \int_{t_{j-1}}^{t_j} S(s) ds = \int_{t_{j-1}}^{t_j} X_j ds = X_j \int_{t_{j-1}}^{t_j} ds = X_j^2.$$

By the renewal reward theorem, we have almost surely that

$$\lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t S(s) ds = \frac{\mathbb{E}[R]}{\mathbb{E}[X]} = \frac{\mathbb{E}[X^2]}{\mathbb{E}[X]}$$

where the fact that $\frac{\mathbb{E}[X^2]}{\mathbb{E}[X]} \geq \mathbb{E}[X]$ yields the inspection paradox.