

Lecture 6: Queueing Theory III

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Disclaimer: These notes are primarily adapted from expository texts, including work by Jyotiprasad Medhi, Karl Sigman, and János Sztrik. These notes are not meant to be complete or fully rigorous; some proofs are not given, incomplete, or only outlined, as they are discussed in class.

Discussion 6.1. Continuing our discussion from Lecture 5 on priorities (Discussion 5.4), now let's consider a non-preemptive case, where we have an M/M/1 queue with two types of customers, but now the arrival of a type 1 customer does not disrupt the service of a type 2 customer. After the server is done serving a type 2 customer, we start serving type 1 customers if there are any. Find the average number of customers and average waiting times of both types.

6.1 Jackson Networks

Jackson's network model is defined as follows. Assume that customers from one node (queueing system) i proceed to an arbitrary node, and new customers may arrive to a node from outside (say customers arrive to node i according to a Poisson process with rate λ_i). Suppose that there are k nodes, where the i th node ($i = 1, \dots, k$) consists of c_i exponential servers with parameter μ_i (that is, each node contains a M/M/ c queueing system). Customers after receiving service at the i th node proceed to the j th node with probability p_{ij} .

Customers at node i depart from the system with probability

$$q_i = 1 - \sum_{j=1}^k p_{ij}.$$

Consider Jackson's general network model with k nodes. The arrivals can be categorized into two groups: the external arrivals (with rate λ_i) and internal arrivals (with rate $\sum_{j=1}^k p_{ji}\lambda_j$). Therefore, the effective arrival rate to node i (or the effective rate of flow through node i) is

$$\alpha_i = \lambda_i + \sum_{j=1}^k p_{ji}\alpha_j, \quad i = 1, 2, \dots, k; \quad (6.1)$$

where these equations are also referred as traffic (flow balance, conservation, etc.) equations.

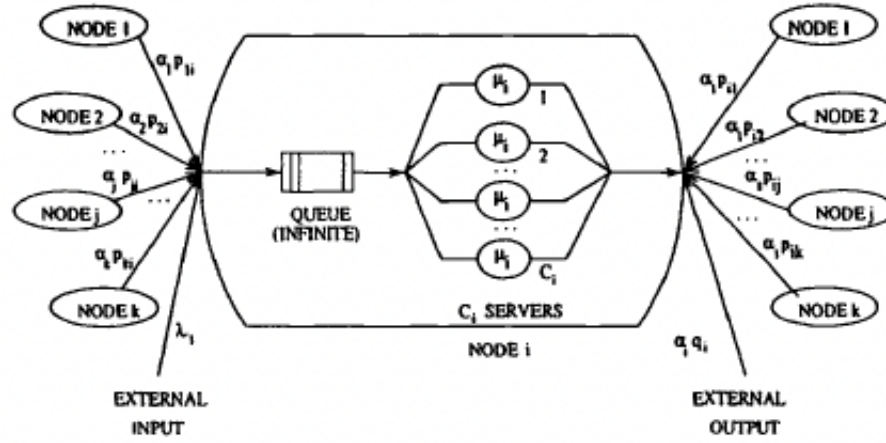


Figure 5.1 Node i in a Jackson Network.

Theorem 6.2 (Jackson's Theorem). Let (n_1, n_2, \dots, n_k) denote the state of the complete system in which there are n_i (in the queue and in service) at node i in a Jackson network of Markovian queues in equilibrium, and let $p(n_1, \dots, n_k)$ be the probability that the system is in the state (n_1, \dots, n_k) .

Assume that

$$\rho_i = \frac{\alpha_i}{\mu_i} < 1, \quad i = 1, 2, \dots, k,$$

where $\{\alpha_i\}$ are given by the balance equations

$$\alpha_i = \lambda_i + \sum_{j=1}^k \alpha_j p_{ji}, \quad i = 1, 2, \dots, k. \quad (6.2)$$

If $p_i(n)$ denotes the probability that there are n customers in the system (in queue plus service) for the $M/M/c_i$ queue with input rate α_i and service rate μ_i for each of the c_i servers, i.e.,

$$p_i(n) = p_i(0) \frac{\left(\frac{\alpha_i}{\mu_i}\right)^n}{n!}, \quad n = 0, 1, 2, \dots, c_i, \quad (6.3)$$

$$= p_i(0) \frac{\left(\frac{\alpha_i}{\mu_i}\right)^n}{c_i! c_i^{n-c_i}}, \quad n = c_i + 1, \dots, \quad (6.4)$$

then we have

$$p(n_1, n_2, \dots, n_k) = p_1(n_1) p_2(n_2) \cdots p_k(n_k). \quad (6.5)$$

Proof Sketch. Let $p_t(n_1, \dots, n_k)$ be the probability that the complete system is in state (n_1, \dots, n_k) at time t .

Let

$$q_i = 1 - \sum_j p_{ij}, \quad a_i(n) = \min\{n_i, c_i\} = \begin{cases} n_i, & \text{if } n < c_i, \\ c_i, & \text{if } n \geq c_i, \end{cases}$$

$$\delta_i = \min\{n_i, 1\} = \begin{cases} 1, & n_i \geq 1, \\ 0, & n_i = 0. \end{cases}$$

Our goal is to write the differential equations satisfied by p_t . Therefore, we will consider the state changes in an infinitesimal interval $(t, t+h)$ following the interval $(0, t)$. p_t . Consider the following four mutually exclusive ways to move from t to $t+h$:

(A) State at t is (n_1, \dots, n_k) and there are no arrivals or departures occur to or from any node externally. We get

$$\Pr(A) = p_t(n_1, \dots, n_k) \left[1 - \left(\sum_i \lambda_i \right) h - \sum_i a_i(n_i) \mu_i h \right] + o(h). \quad (6.6)$$

(B) State at t is $(n_1, \dots, n_i + 1, \dots, n_k)$ and there is one service completion at i in $(t, t+h)$, and this completion departs from the system (with probability q_i). We get

$$\Pr(B) = \sum_{i=1}^k p_t(n_1, \dots, n_i + 1, \dots, n_k) [a_i(n_i + 1) \mu_i q_i] h + o(h). \quad (6.7)$$

(C) State at t is $(n_1, \dots, n_i - 1, \dots, n_k)$ and there is one arrival from the external source to node i in the interval $(t, t+h)$. We get

$$\Pr(C) = \sum_{i=1}^k p_t(n_1, \dots, n_i - 1, \dots, n_k) [\lambda_i h \delta_i] + o(h). \quad (6.8)$$

(D) State at t is $(n_1, \dots, n_i + 1, \dots, n_j - 1, \dots, n_k)$: there is one service completion at node i in $(t, t+h)$, and the one whose service is completed moves to node j with probability p_{ij} . Thus,

$$\Pr(D) = \sum_i \sum_j p_t(n_1, \dots, n_i + 1, \dots, n_j - 1, \dots, n_k) [a_i(n_i + 1) \mu_i h p_{ij}] + o(h). \quad (6.9)$$

Merging all cases, we have

$$p_{t+h}(n_1, \dots, n_k) = \Pr(A) + \Pr(B) + \Pr(C) + \Pr(D). \quad (6.10)$$

(6.10) can be written as

$$p_{t+h}(n_1, \dots, n_k) - \Pr(A) = \Pr(B) + \Pr(C) + \Pr(D).$$

(6.11)

$$p_{t+h}(n_1, \dots, n_k) - p_t(n_1, \dots, n_k) \left[1 - \left(\sum_i \lambda_i \right) h - \sum_i a_i(n_i) \mu_i h \right] - o(h) = \Pr(B) + \Pr(C) + \Pr(D).$$

(6.12)

Taking the limit as $h \rightarrow 0$ and solving for $p'(t) = 0$ gives the equations satisfied by steady-state probabilities:

$$\begin{aligned} \left[\sum_i \lambda_i + \sum_i a_i(n_i) \mu_i \right] p(n_1, \dots, n_k) &= \sum_i a_i(n_i + 1) \mu_i q_i p(n_1, \dots, n_i + 1, \dots, n_k) \\ &+ \sum_i \lambda_i \delta_i p(n_1, \dots, n_i - 1, \dots, n_k) \\ &+ \sum_i \sum_j a_i(n_i + 1) \mu_i p_{ij} p(n_1, \dots, n_i + 1, \dots, n_j - 1, \dots, n_k). \end{aligned}$$

Finally, one can show that (6.5) uniquely satisfies the equations above.

■

6.2 The M/G/1 Model

Now assume that we have a Poisson arrival process with rate λ , the service times are i.i.d. and follow a general distribution with mean $\mathbb{E}[S] = \frac{1}{\mu}$, and there is a single server. Again, for stability, we assume that $\rho = \frac{\lambda}{\mu} < 1$.

Let R be the residual service time and let P_k denote the probability that there are k customers in the system in the steady-state. By PASTA property, we have

$$\begin{aligned} W_q &= \sum_{k=1}^{\infty} (E(R) + (k-1)E(S)) P_k \\ &= \sum_{k=1}^{\infty} E(R) P_k + \left(\sum_{k=1}^{\infty} (k-1) P_k \right) E(S) \\ &= E(R) \rho + L_q E(S). \end{aligned}$$

where the equation follows since $1 - \rho = P_0$. By Little's Law, we get

$$W_q = \frac{\rho \mathbb{E}[R]}{1 - \rho} \quad (6.13)$$

which is known as the Pollaczek-Khintchine mean value formula. So now we have to characterize $\mathbb{E}[R]$.

Proposition 6.3. *We have $\mathbb{E}[R] = \frac{\mathbb{E}[S^2]}{2\mathbb{E}[S]}$.*

There are various ways to prove this, but the most pedagogical one is via renewal reward theorem (see the next section). From Proposition 6.3, we have

$$E(R) = \frac{E(S^2)}{2E(S)} = \frac{\text{Var}(S) + E[S]^2}{2E(S)} = \frac{1}{2} (C_S^2 + 1) E(S), \quad (2.33)$$

where C_S^2 is the squared coefficient of the service time S . From here, we get

Thus for the mean waiting time we have

$$W_q = \frac{\rho E(R)}{1 - \rho} = \frac{\rho}{2(1 - \rho)} (C_S^2 + 1) E(S).$$

And by Little's law, we get

$$L_q = \frac{\rho^2}{1 - \rho} \frac{C_S^2 + 1}{2}.$$

Discussion 6.4. *Kingman's G/G/N formula:*

$$W \approx \frac{1}{\mu N} \cdot \frac{\rho \sqrt{2(N+1)} - 1}{1 - \rho} \cdot \frac{C_A^2 + C_S^2}{2} + \frac{1}{\mu}.$$

Discussion 6.5. *Matching queues, and the relationship between regret and queue-lengths.*

6.3 The Renewal Reward Theorem

Definition 6.6. *A random point process $\psi = \{t_n\}$ for which the (non-negative) interarrival times $X_n = t_n - t_{n-1}$, $n \geq 1$, form an i.i.d. sequence is called a renewal process.*

Following the definition, t_n is called the n th *renewal epoch* and $F(x) := P(X \leq x)$, $x \geq 0$, denotes the common inter-arrival time distribution. $t_n = X_1 + \dots + X_n$, and $N(t) = \max\{n : t_n \leq t\}$ is the counting process. The rate of the renewal process is denoted by $\lambda \triangleq 1/\mathbb{E}[X]$.

Theorem 6.7 (Elementary renewal theorem). *For a renewal process,*

$$\lim_{t \rightarrow \infty} \frac{N(t)}{t} = \lambda \quad a.s.$$

and

$$\lim_{t \rightarrow \infty} \frac{\mathbb{E}[N(t)]}{t} = \lambda.$$

Proof. We start with the first statement. Note that $t_n = X_1 + \cdots + X_n$, $n \geq 1$. Consider the time

$$t_{N(t)} \leq t < t_{N(t)+1}, \quad (6.14)$$

Then we have

$$t_{N(t)} = \sum_{j=1}^{N(t)} X_j, \quad t_{N(t)+1} = \sum_{j=1}^{N(t)+1} X_j,$$

which can be written as

$$\frac{1}{N(t)} \sum_{j=1}^{N(t)} X_j \leq \frac{t}{N(t)} < \frac{1}{N(t)} \sum_{j=1}^{N(t)+1} X_j.$$

By the strong law of large numbers, the left and right hand sides converge to $E(X)$ as $t \rightarrow \infty$ almost surely. One can prove the second part via Wald's identity (which you will prove in your homework with some provided hints). ■

Now let $R(t) = \sum_{j=1}^{N(t)} R_j$ be the total amount of reward collected by time t , where $N(t)$ is the counting process for the renewal process. We want to calculate our long-run reward rate

$$\lim_{t \rightarrow \infty} \frac{R(t)}{t}.$$

Theorem 6.8 (Renewal reward theorem). *For a positive recurrent renewal process in which a reward R_j is earned during cycle length X_j and such that $\{(X_j, R_j) : j \geq 1\}$ is i.i.d. with $\mathbb{E}[|R_j|] < \infty$, the long run rate at which rewards are earned is given by*

$$\lim_{t \rightarrow \infty} \frac{R(t)}{t} = \frac{\mathbb{E}[R]}{\mathbb{E}[X]} = \lambda \mathbb{E}[R] \quad a.s., \quad (6.15)$$

where (X, R) denotes a typical “cycle” (X_j, R_j) ; $\lambda = \{\mathbb{E}[X]\}^{-1}$ is the arrival rate for the renewal process.

Moreover,

$$\lim_{t \rightarrow \infty} \frac{\mathbb{E}[R(t)]}{t} = \frac{\mathbb{E}[R]}{\mathbb{E}[X]}. \quad (6.16)$$

Discussion 6.9. Now let's prove Proposition 6.3 via the renewal reward theorem. Consider a renewal point process $\{t_n : n \geq 1\}$ with i.i.d. interarrival times $X_n = t_n - t_{n-1}$, $n \geq 1$. Define

$$A(t) = t_{N(t)+1} - t, \quad t \geq 0. \quad (6.17)$$

$A(t)$ is called the excess at time t , or remaining lifetime. If $t_{n-1} \leq t < t_n$, then

$$A(t) = t_n - t \leq X_n.$$

Note that if $\{t_n\}$ is a Poisson process at rate λ , then by the memoryless property we have $A(t) \sim \exp(\lambda)$, $t \geq 0$. But for a general renewal process (as in residual service time in $M/G/1$ queue), we need to be smarter. We want to show that

$$\lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t A(s) ds = \frac{\mathbb{E}[X^2]}{2\mathbb{E}[X]} \quad \text{a.s.}$$

Note that we can view the i.i.d. X_j as cycle lengths (service times), and $r(t) = A(t)$ as the generated reward rate at time t . Let R_1 be the generated reward in the first cycle. Then we have

$$R_1 = \int_0^{X_1} A(s) ds = \int_0^{X_1} (X_1 - s) ds = \frac{X_1^2}{2}.$$

Since $\{(X_j, R_j)\}$'s are i.i.d., by the renewal reward theorem, we almost surely have

$$\lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t A(s) ds = \frac{\mathbb{E}[R]}{\mathbb{E}[X]} = \frac{\mathbb{E}[X^2]}{2\mathbb{E}[X]}$$

Discussion 6.10 (Inspection paradox). Let $S(t) = t_{N(t)+1} - t_{N(t)}$ be the length of the interarrival time covering time t . If $t_{j-1} \leq t < t_j$, then we have $S(t) = X_j$. Define the reward rate as $r(t) = S(t)$. Then we get

$$R_j = \int_{t_{j-1}}^{t_j} S(s) ds = \int_{t_{j-1}}^{t_j} X_j ds = X_j \int_{t_{j-1}}^{t_j} ds = X_j^2.$$

By the renewal reward theorem, we have almost surely that

$$\lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t S(s) ds = \frac{\mathbb{E}[R]}{\mathbb{E}[X]} = \frac{\mathbb{E}[X^2]}{\mathbb{E}[X]}$$

where the fact that $\frac{\mathbb{E}[X^2]}{\mathbb{E}[X]} \geq \mathbb{E}[X]$ yields the inspection paradox.